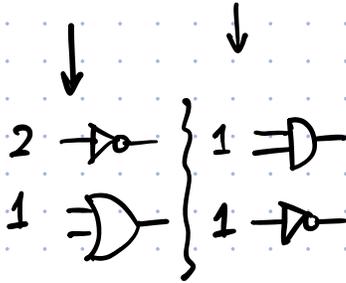


February 11, 2026

Focus: Digital logic optimization

|   |   |         |         |
|---|---|---------|---------|
| a | b | $a'+b'$ | $(ab)'$ |
| 0 | 0 | 1       | 1       |
| 0 | 1 | 1       | 1       |
| 1 | 0 | 1       | 1       |
| 1 | 1 | 0       | 0       |



\* Boolean Algebra - Axioms

|     | AND Form   | OR Form   |   |
|-----|--|---|---|
| (i) | $1x = x$<br>$0x = \emptyset$<br>$xx = x$<br>$x x' = \emptyset$ | $0+x = x$<br>$1+x = 1$<br>$x+x = x$<br>$x+x' = 1$ | Identity Law<br>Null Law<br>Idempotent Law<br>Inverse Law |

(ii)  $x'' = x$  Double complement law

| (iii) | AND Form  | OR Form  |  |
|-------|---|--|--|
|       | $xy = yx$<br>$(xy)z = x(yz)$<br>$x+yz = (x+y)(x+z)$ | $x+y = y+x$<br>$(x+y)+z = x+(y+z)$<br>$x(y+z) = xy+xz$ | commutative law<br>associative law<br>distributive law |

|      |                                 |                               |                                  |
|------|---------------------------------|-------------------------------|----------------------------------|
| (iv) | $x(x+y) = x$<br>$(xy)' = x'+y'$ | $x+xy = x$<br>$(x+y)' = x'y'$ | Absorption Law<br>DeMorgan's Law |
|------|---------------------------------|-------------------------------|----------------------------------|

\* Example:  $f(x, y, z) = xy + x'z + yz \rightsquigarrow 5 \text{ gates}$

$$\begin{aligned}
 &= xy + x'z + yz \cdot 1 && \text{identity law} \\
 &= xy + x'z + yz(x+x') && \text{inverse law} \\
 &= xy + x'z + yzx + yzx' && \text{distributive law} \\
 &= xy + x'z + xyx + x'yx && \text{commutative law} \\
 &= xy + x'z + (xy)x + (x'z)y && \text{associative law} \\
 &= xy + (xy)x + x'z + (x'z)y && \text{commutative} \\
 &= xy(1+z) + x'z(1+y) && \text{distributive} \\
 &= xy + x'z \rightsquigarrow 4 \text{ gates} && \text{null } \times 2
 \end{aligned}$$

\* Example:  $f(x, y, z) = (x+y)(x+z)$

$$\begin{aligned}
 &= (x+y)x \\
 &= x+xy \\
 &= x
 \end{aligned}$$

\*  $f(x, y, z) = xy'z' + xy'z + xyz' + xyz$

$$\begin{aligned}
 &= x(y'z' + y'z + yz' + yz) \\
 &= x(y'(z'+z) + y(z'+z)) \\
 &= x(y' \cdot 1 + y \cdot 1) \\
 &= x1 \\
 &= x
 \end{aligned}$$

\* Algebraic approach  $\rightarrow$  not obvious what to do next.  
 $\rightarrow$  systematic approach

Boolean expression simplification has been known to be NP-complete.